



Algoritmi e Strutture Dati

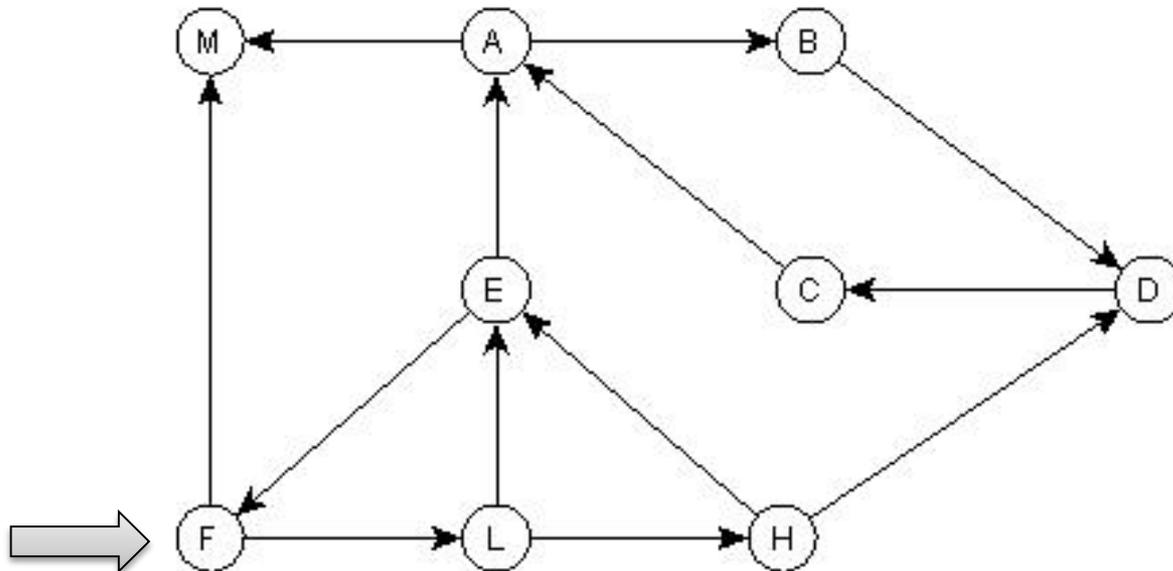
Esercitazione 8

Domenico Fabio Savo

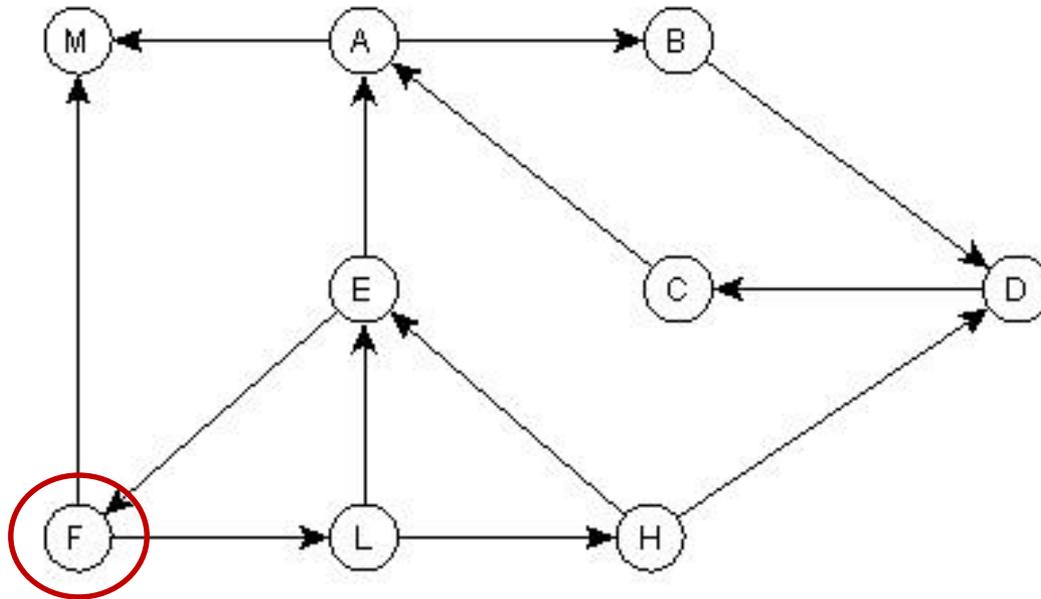
Esercizio: visita in ampiezza

Esercizio:

- Eseguire la **visita in ampiezza** del grafo *orientato* rappresentato in figura partendo dal nodo **F**.
- Presentare l'**albero BFS** prodotto dalla visita.



Esercizio: visita in ampiezza

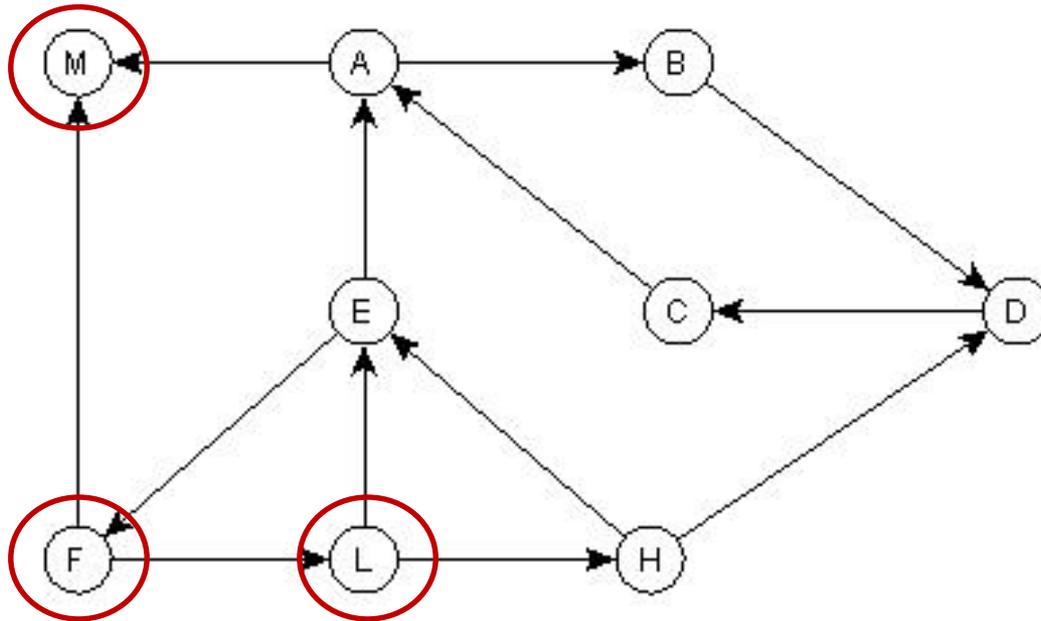


Albero BST

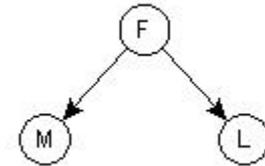
F

F

Esercizio: visita in ampiezza

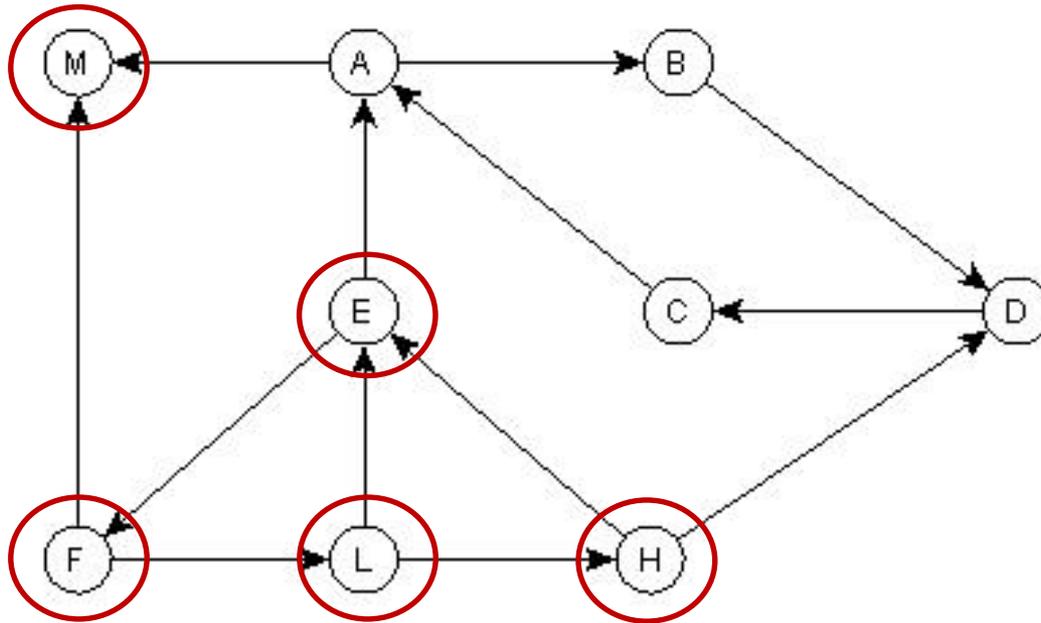


Albero BST

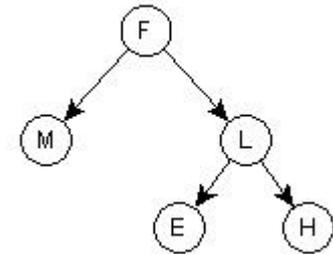


M L F

Esercizio: visita in ampiezza

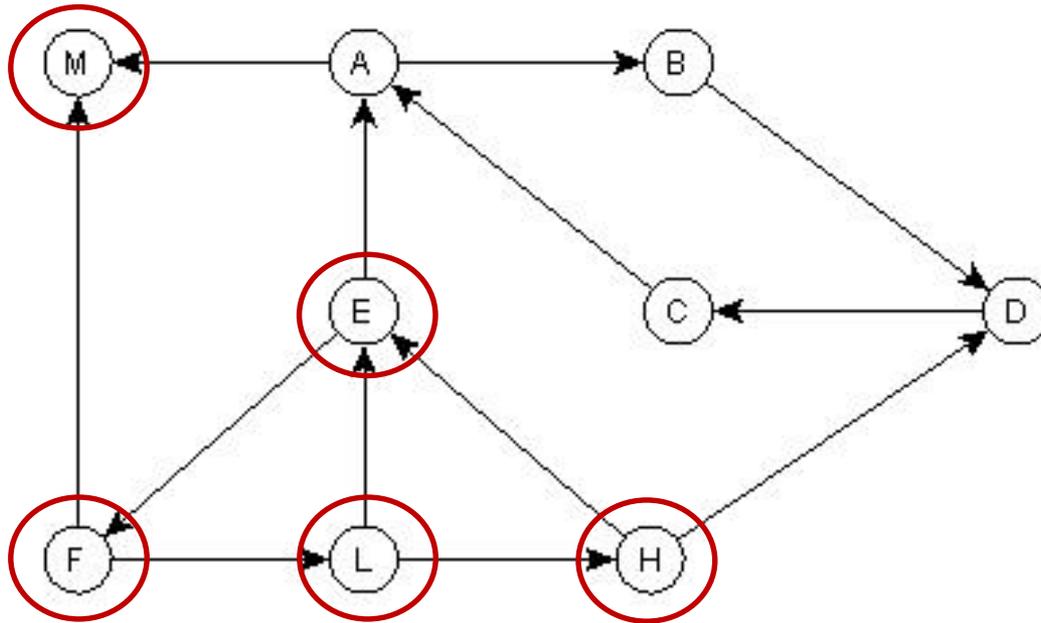


Albero BST

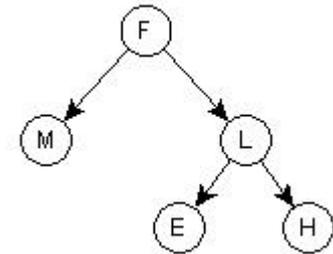


H E M L F

Esercizio: visita in ampiezza

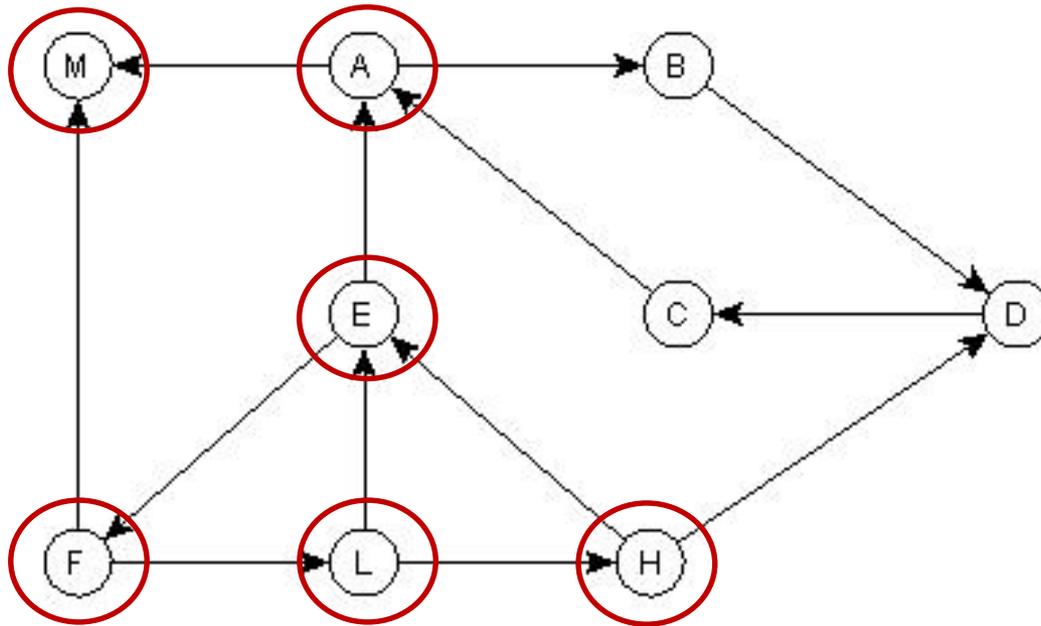


Albero BST

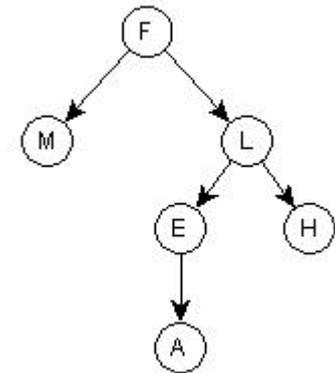


H E M L F

Esercizio: visita in ampiezza

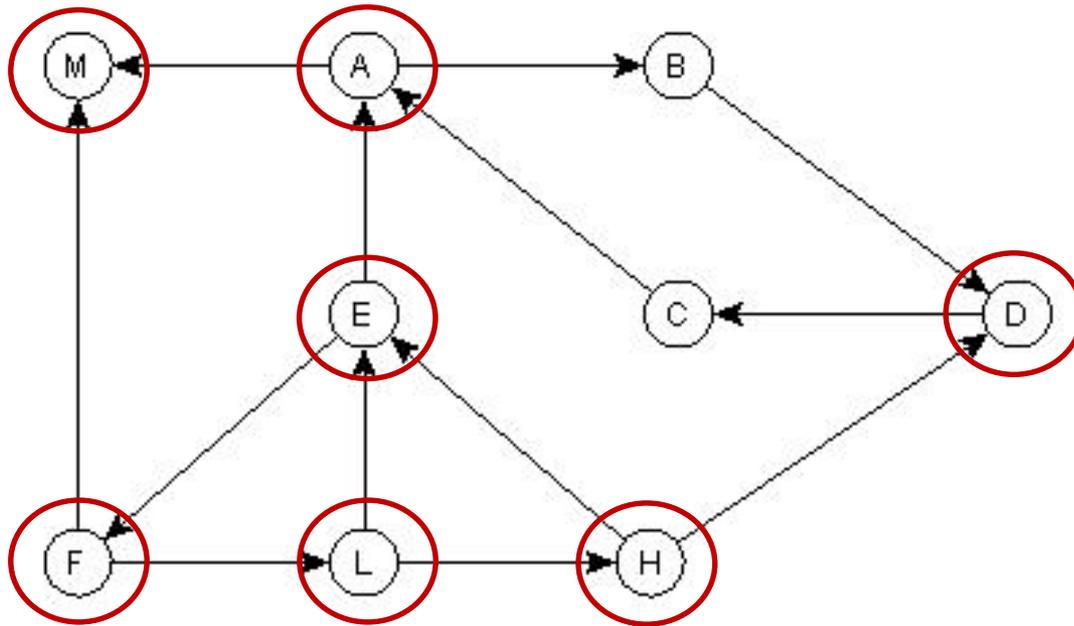


Albero BST

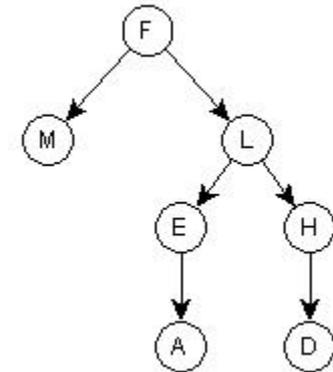


A H E M L F

Esercizio: visita in ampiezza

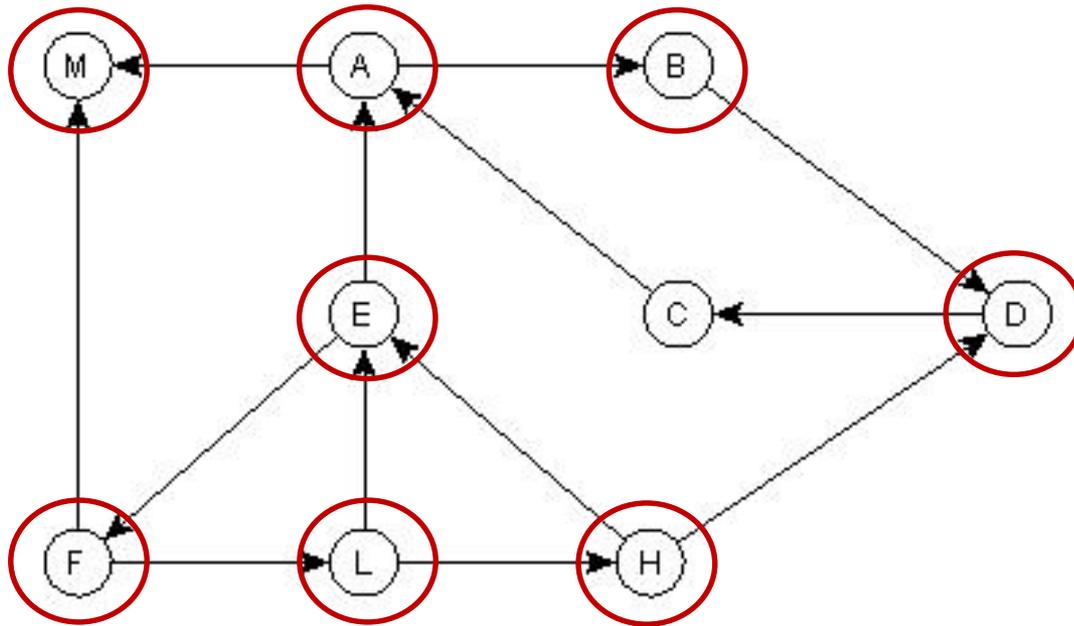


Albero BST

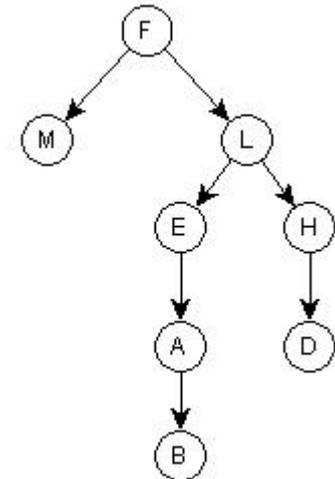


DAHEMLF

Esercizio: visita in ampiezza

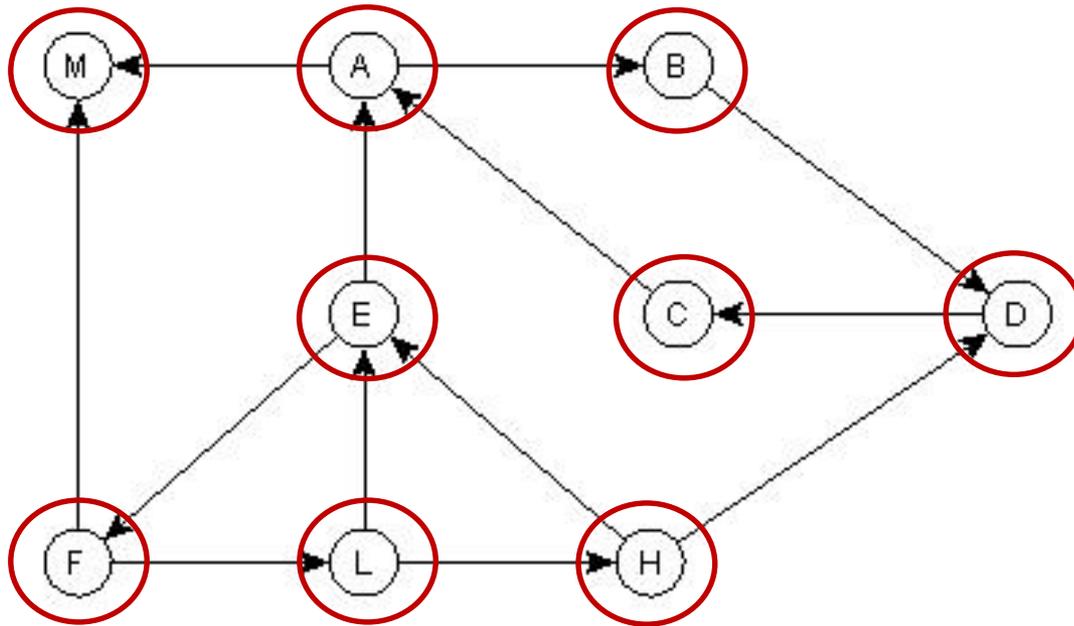


Albero BST

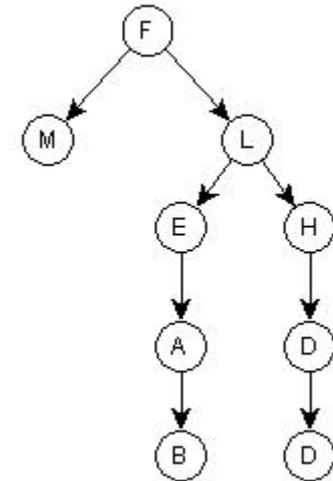


B D A H E M L F

Esercizio: visita in ampiezza

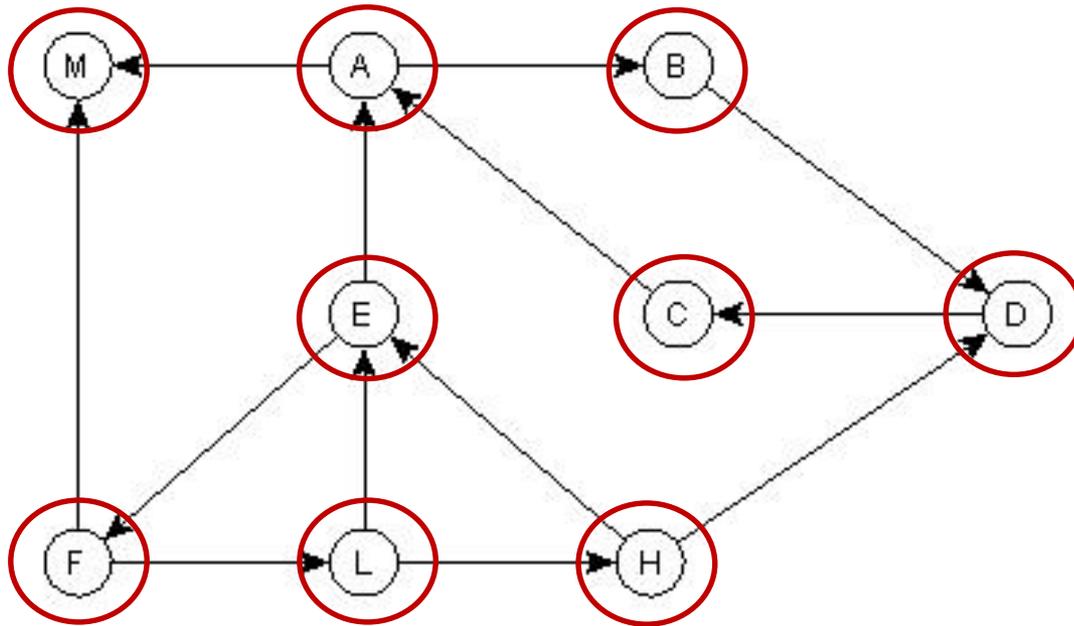


Albero BST

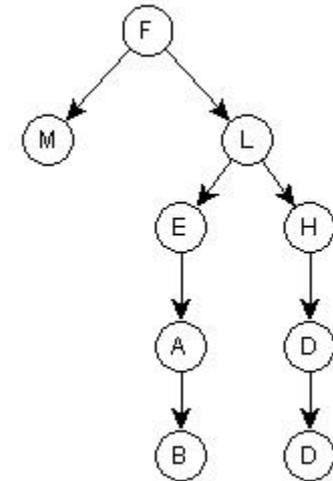


C B D A H E M L F

Esercizio: visita in ampiezza

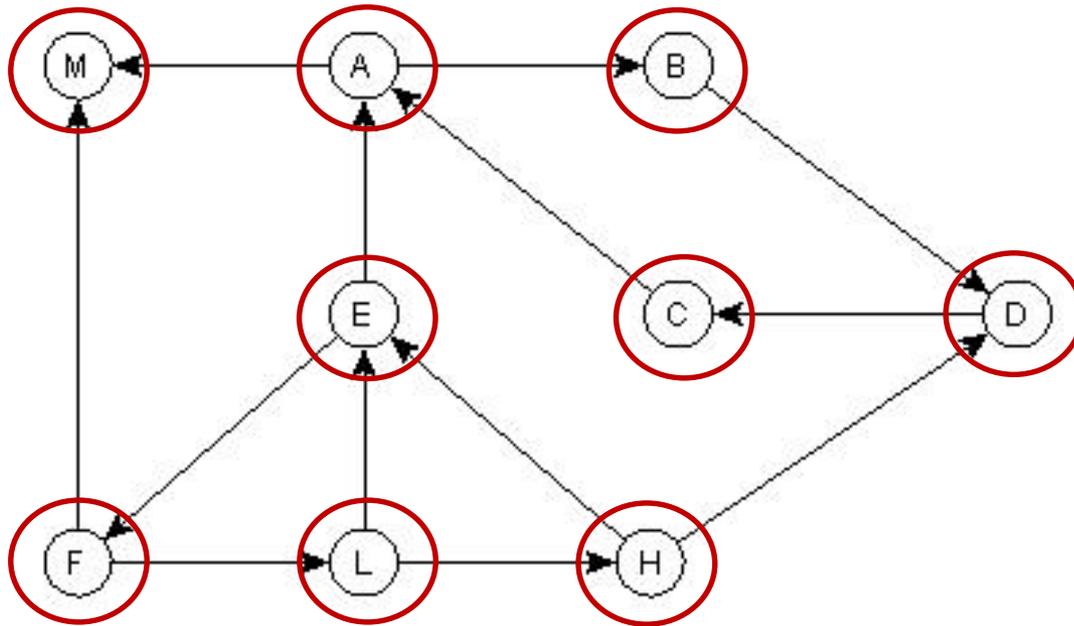


Albero BST



C B D A H E M L F

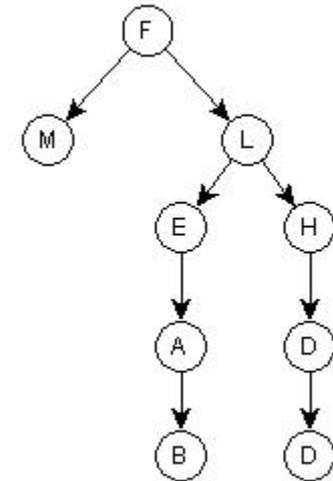
Esercizio: visita in ampiezza



C B D A H E M L F

La coda è vuota!

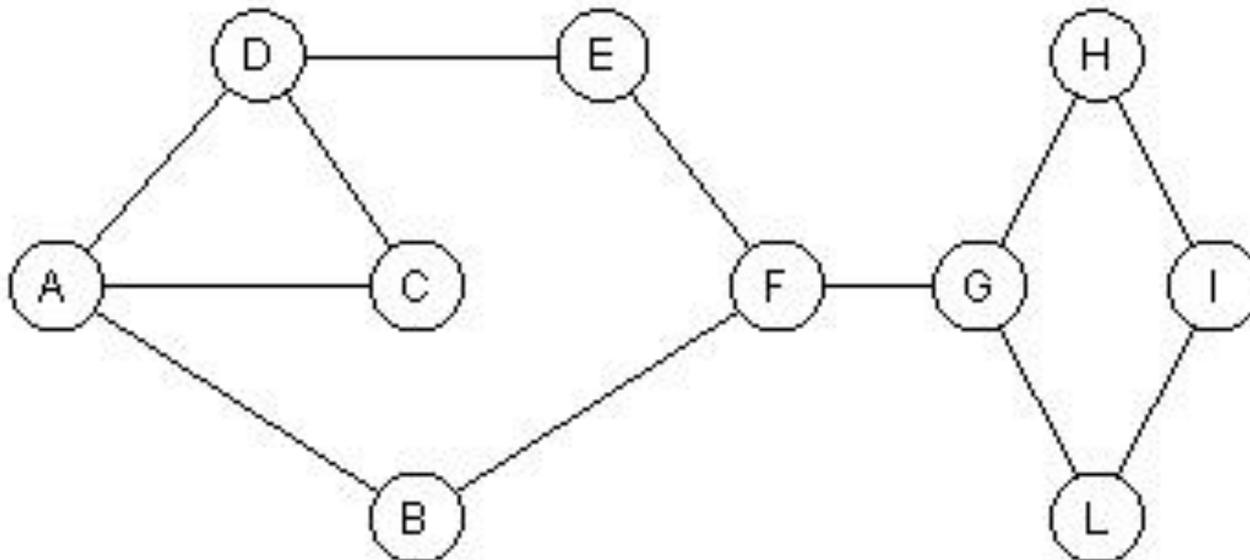
Albero BST



Esercizio: visita in profondità

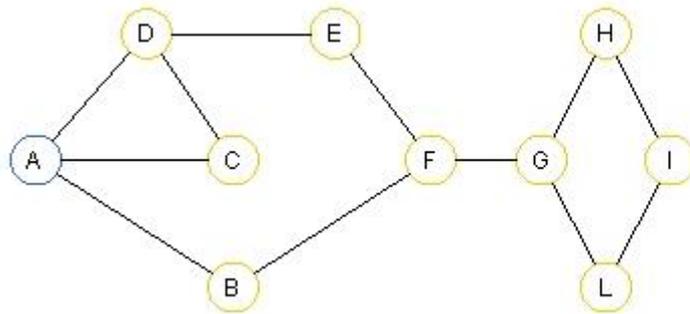
Esercizio:

- Eseguire la **visita in profondità** del grafo *non orientato* rappresentato in figura partendo dal nodo **A**.
- Presentare l'**albero DFS** prodotto dalla visita.



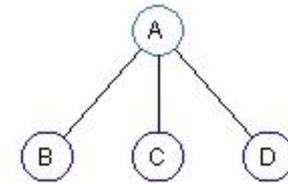
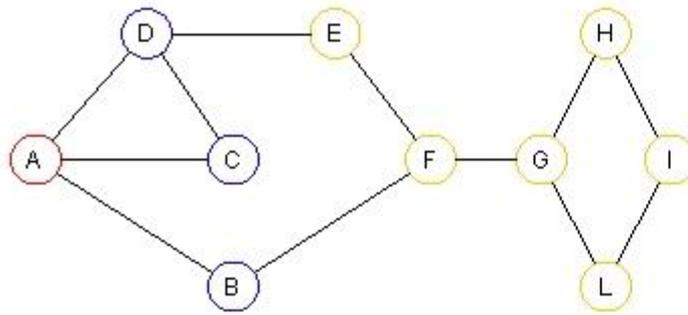


Esercizio: visita in profondità





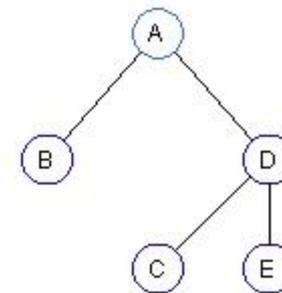
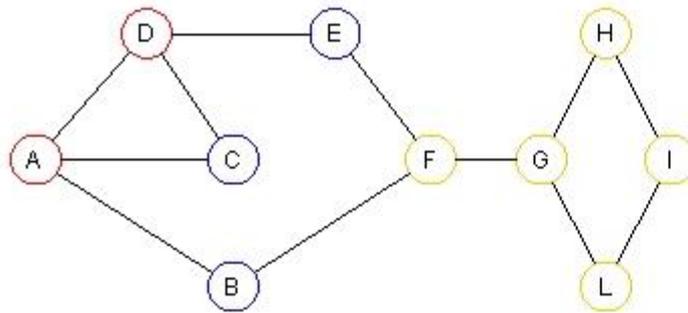
Esercizio: visita in profondità



D
C
B
A



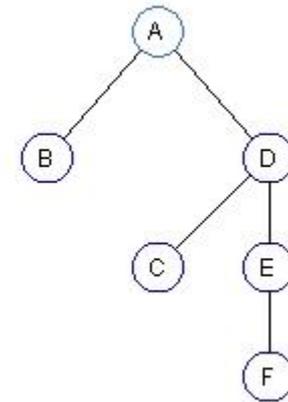
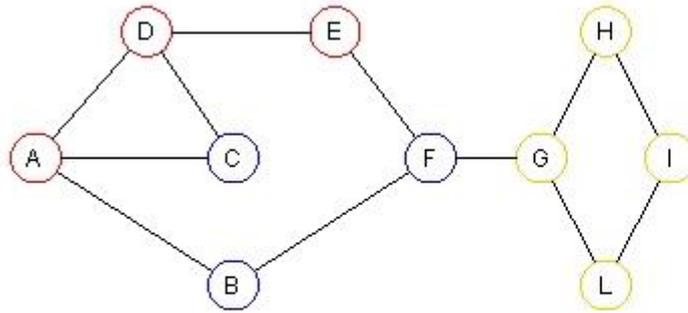
Esercizio: visita in profondità



E
C
D
C
B
A

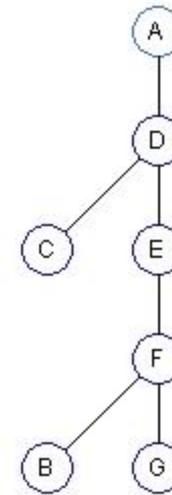
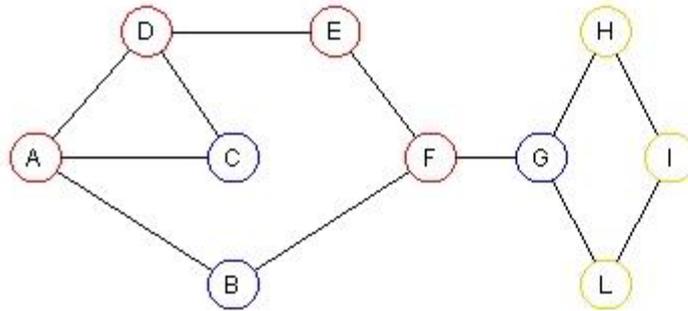


Esercizio: visita in profondità



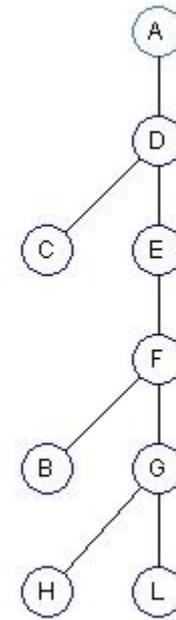
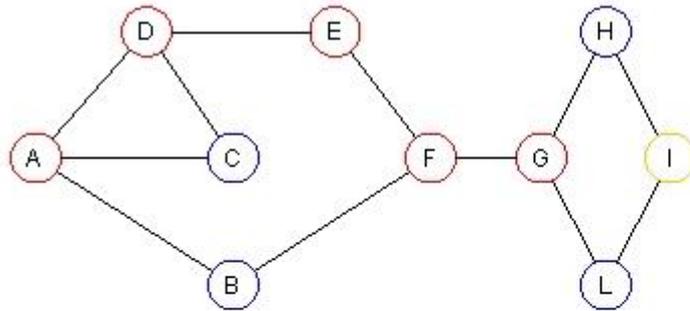
F
E
C
D
C
B
A

Esercizio: visita in profondità



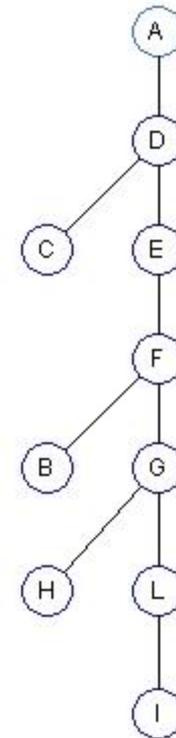
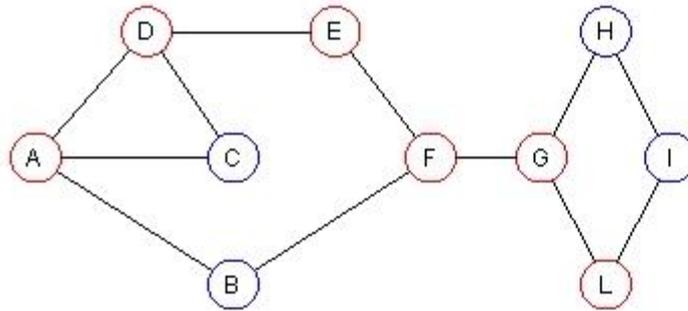
G
B
F
E
C
D
C
B
A

Esercizio: visita in profondità



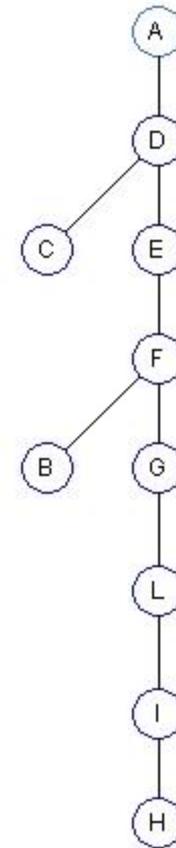
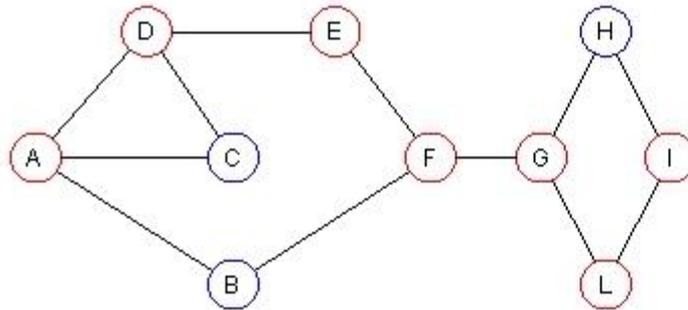
L
H
G
B
F
E
C
D
C
B
A

Esercizio: visita in profondità



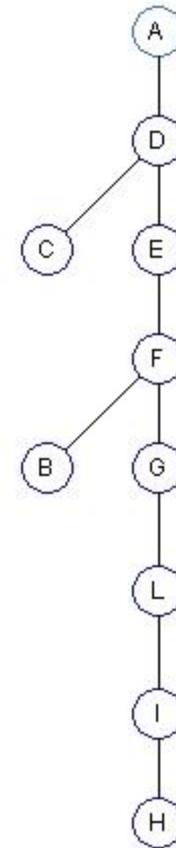
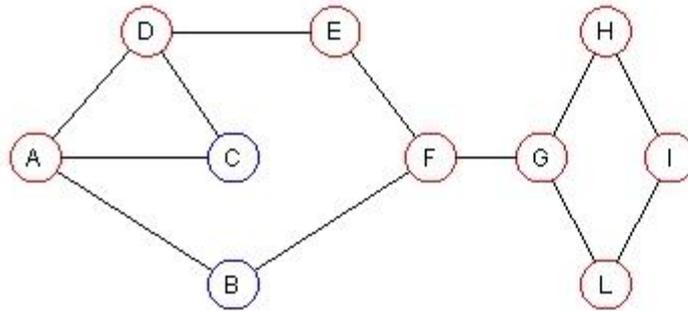
I
L
H
G
B
F
E
C
D
C
B
A

Esercizio: visita in profondità



H
I
L
H
G
B
F
E
C
D
C
B
A

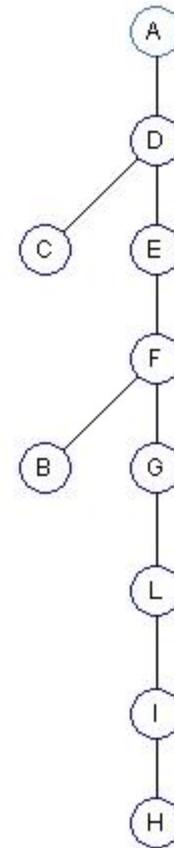
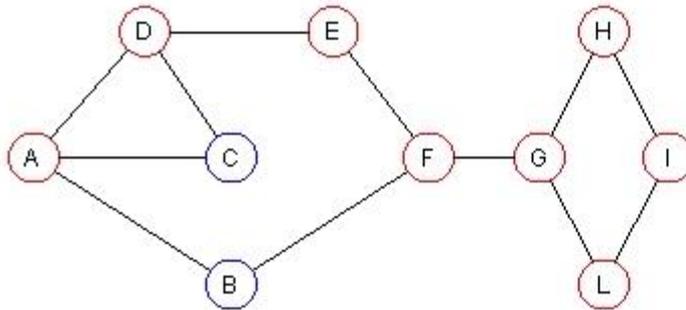
Esercizio: visita in profondità



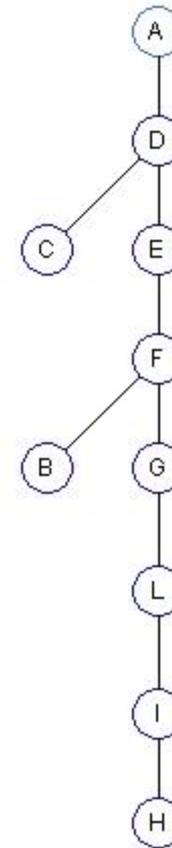
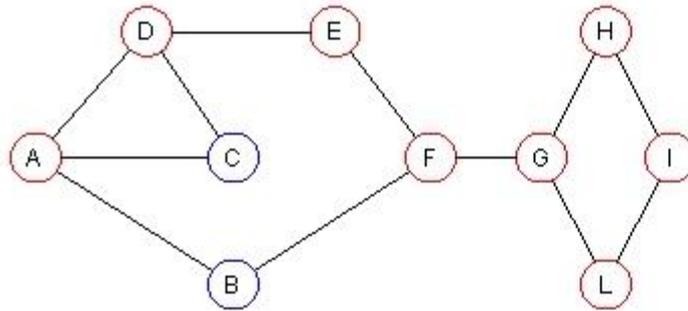
H
I
L
H
G
B
F
E
C
D
C
B
A

Esercizio: visita in profondità

H
I
L
H
G
B
F
E
C
D
C
B
A

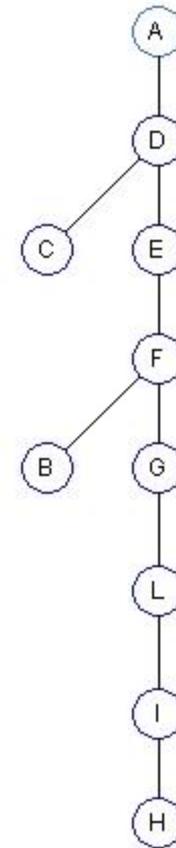
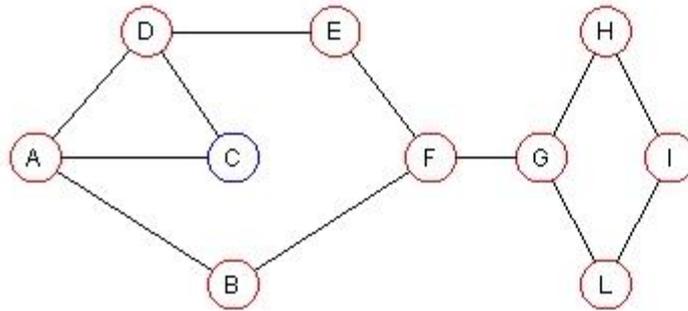


Esercizio: visita in profondità



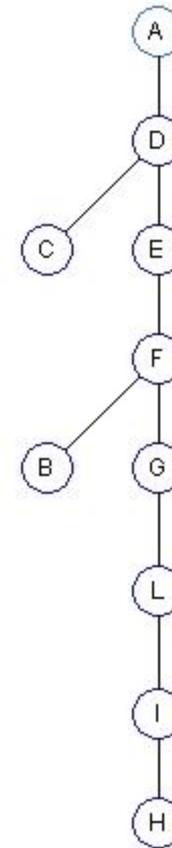
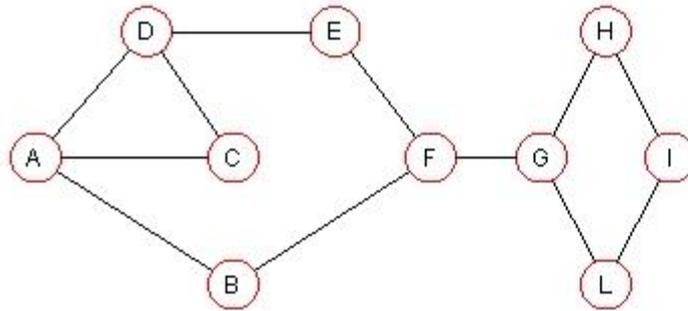
H
I
L
H
G
B
F
E
C
D
C
B
A

Esercizio: visita in profondità



H
I
L
H
G
B
F
E
C
D
C
B
A

Esercizio: visita in profondità

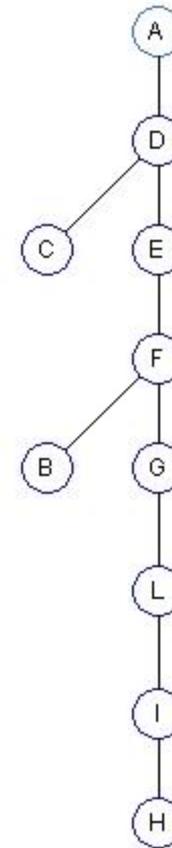
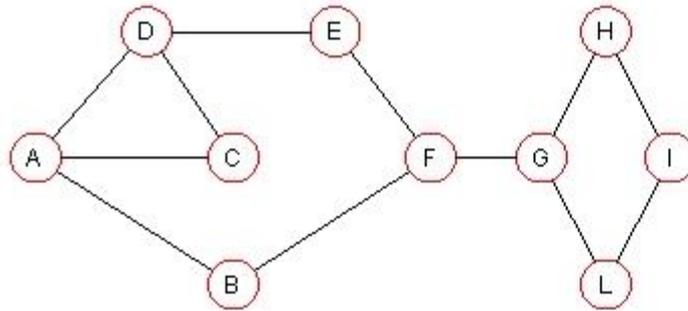


H
I
L
H
G
B
F
E
C
D
C
B
A

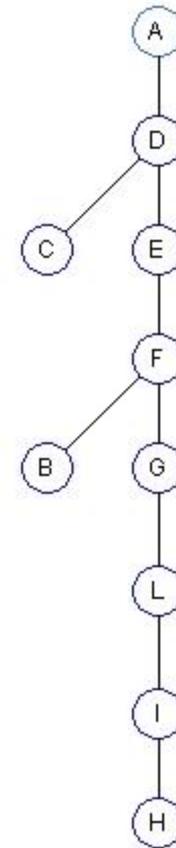
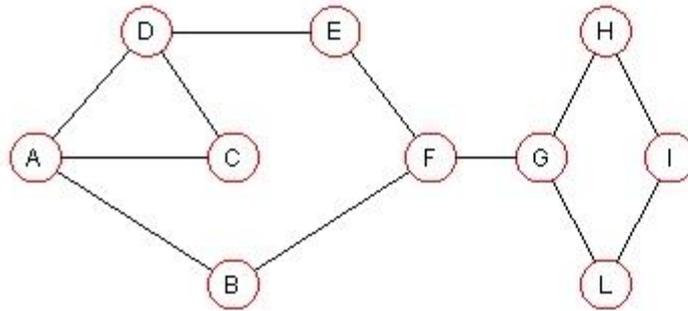


Esercizio: visita in profondità

H
I
L
H
G
B
F
E
C
D
C
B
A



Esercizio: visita in profondità

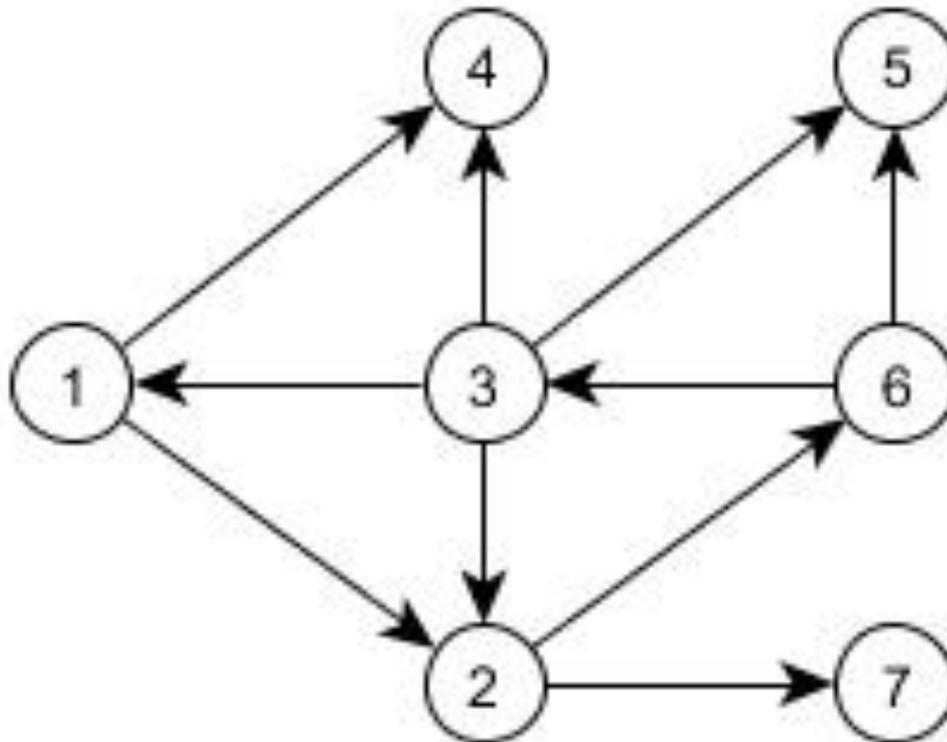


H
I
L
H
G
B
F
E
C
D
C
B
A

La pila è vuota!

Esercizio (max grado)

Sia $G = (V,E)$ il grafo *orientato* rappresentato in figura.



Esercizio (max grado)

Si chiede di:

1. Fornire una rappresentazione del grafo basata sulla matrice di adiacenza.
2. Scrivere l'algoritmo
nodoMaxGrado (array G[n][n]) → intero v
che, preso in ingresso un grafo G (rappresentato come una matrice di adiacenza), restituisce l'indice (indicato come un intero) di uno dei vertici con *grado in uscita massimo*.
3. Indicare il costo computazionale dell'algoritmo proposto.

Nota: un vertice **v** ha *grado in uscita massimo* se **non** esiste in **G** nessun vertice **w** tale che $\delta_{\text{out}}(\mathbf{v}) < \delta_{\text{out}}(\mathbf{w})$

Esercizio (max grado)

Soluzione:

1. Fornire una rappresentazione del grafo con **matrice di adiacenza**.

	1	2	3	4	5	6	7
1	0	1	0	1	0	0	0
2	0	0	0	0	0	1	1
3	1	1	0	1	1	0	0
4	0	0	0	0	0	0	0
5	0	0	0	0	0	0	0
6	0	0	1	0	1	0	0
7	0	0	0	0	0	0	0

Esercizio (max grado)

Soluzione:

```
algoritmo nodoMaxGrado (array  $G[n][n]$ )  $\rightarrow$  intero  $v$ 
1. intero  $v\_max \leftarrow 1$ ;
2. intero  $g\_max \leftarrow 0$ ;
3. for (int  $r$ ;  $r = 1$ ;  $r \leq n$ ) do
4.     int  $current\_g \leftarrow 0$ ;
5.     for(int  $c$ ;  $c = 1$  ;  $c \leq n$  ) do
6.          $current\_g \leftarrow current\_g + G[r][c]$ ;
7.     if(  $current\_g > g\_max$ ) then
8.          $g\_max \leftarrow current\_g$ ;
9.          $v\_max \leftarrow r$ ;
10. return  $v\_max$ ;
```

\rightarrow Il costo dell'algoritmo è $O(n^2)$