

Turing Machines

- In 1936, A.M. Turing proposed the Turing machine as a model of *any possible computation*.
- This model was *built* using electro-mechanical devices after several years.
- Turing machine long has been recognized as an accurate model for what any physical computing device is capable of doing.

Turing Machines

Turing formalized the idea of “mechanical procedure”:

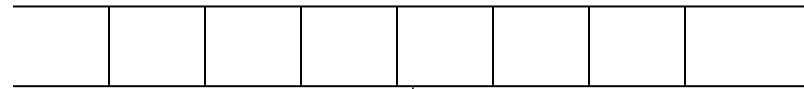
- Can be described by a finite number of instructions.
- Instructions are simple and mechanical.
- Have a finite number of internal states.
- Can deal with input not restricted in size.
- Have an unlimited storage space for calculations.
- Can produce output of unlimited size.

Turing Machine

- Not a real machine, but a model of computation
- Components:
 - 1-way infinite tape: unlimited memory
 - Store input, output, and intermediate results
 - Infinite cells
 - Each cell has a symbol from a finite alphabet
- Tape head:
 - Point to one cell
 - Read or write a symbol to that cell
 - move left or right

Turing Machines

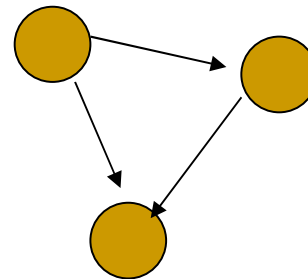
This tape is for input, storage and output



Tape head



Finite
Control



Definition: A Turing Machine is a 7-tuple

$T = (Q, \Sigma, \Gamma, \delta, q_0, q_{\text{accept}}, q_{\text{reject}})$, where:

Q is a finite set of states

(Σ is the input alphabet, where $\square \notin \Sigma$ (\square blank))

Γ is the tape alphabet, where $\square \in \Gamma$ and $\Sigma \subseteq \Gamma$

$\delta : Q \times \Gamma \rightarrow Q \times \Gamma \times \{L,R\}$

$q_0 \in Q$ is the start state

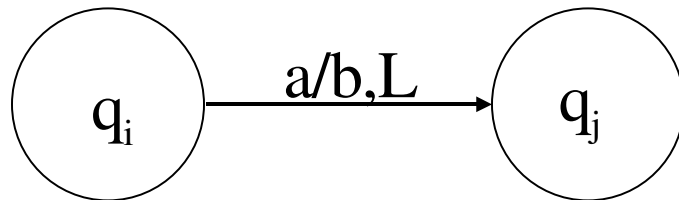
$q_{\text{accept}} \in Q$ is the accept state

$q_{\text{reject}} \in Q$ is the reject state, and $q_{\text{reject}} \neq q_{\text{accept}}$

Turing Machines

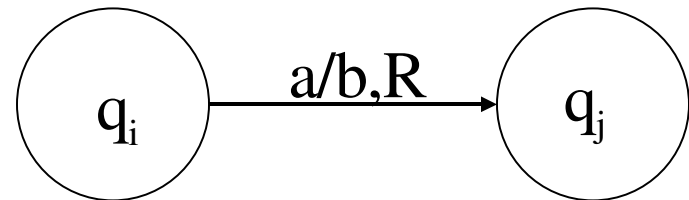
δ is a **move function** mapping from $Q \times \Gamma$ to $Q \times \Gamma \times \{L, R\}$

Replace the current symbol “a” by “b”,
and move to the left.



$$\delta(q_i, a) = (q_j, b, L)$$

Replace the current symbol “a” by “b”,
and move to the right.



$$\delta(q_i, a) = (q_j, b, R)$$

Note that if $a = b$, we write “a,L” instead of “a/a,L” along the edge.

A TM for **Successor** Program

Sample Rules:

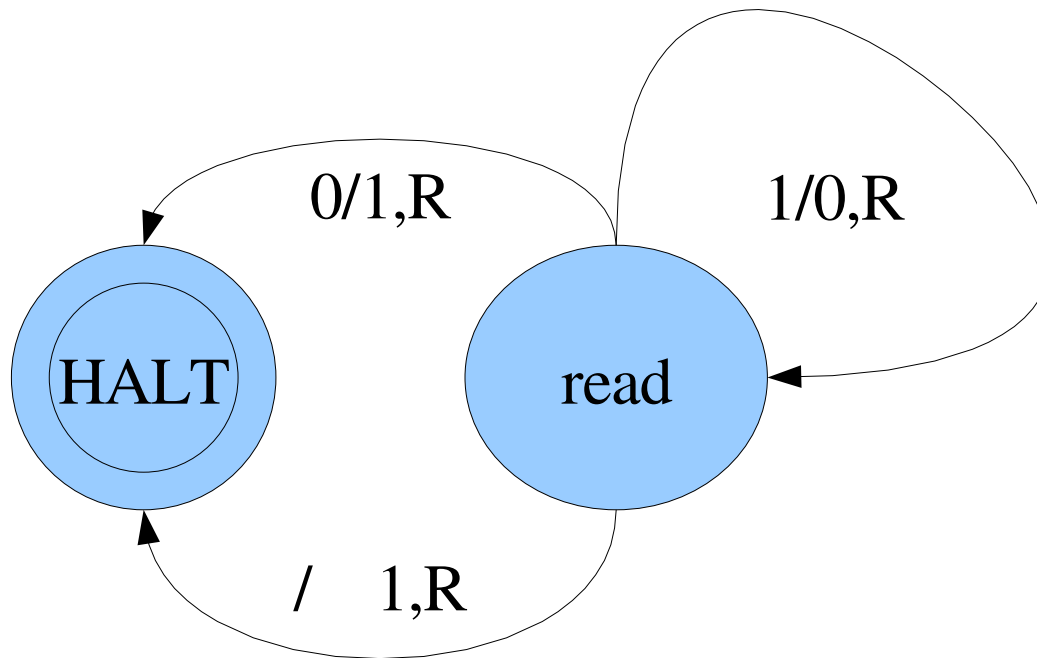
If read 1, write 0, go right, repeat.

If read 0, write 1, HALT!

If read \square , write 1, HALT!

Let's see how they are carried out on a piece of paper that contains the *reverse* binary representation of 47:

TM successor



If read 1, write 0, go right, repeat.

If read 0, write 1, HALT!

If read , write 1, HALT!



If read 1, write 0, go right, repeat.

If read 0, write 1, HALT!

If read , write 1, HALT!



If read 1, write 0, go right, repeat.

If read 0, write 1, HALT!

If read , write 1, HALT!



If read 1, write 0, go right, repeat.

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If read 1, write 0, go right, repeat.

If read 0, write 1, HALT!

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If read 1, write 0, go right, repeat.

If read 0, write 1, HALT!

If read , write 1, HALT!



So the successor's output on 111101 was 000011 which is the reverse binary representation of 48.

Similarly, the successor of 127 should be 128:

If read 1, write 0, go right, repeat.

If read 0, write 1, HALT!

If read \square , write 1, HALT!



If read 1, write 0, go right, repeat.

If read 0, write 1, HALT!

If read , write 1, HALT!



If read 1, write 0, go right, repeat.

If read 0, write 1, HALT!

If read , write 1, HALT!

0	0	1	1	1	1	1			
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If read 1, write 0, go right, repeat.

If read 0, write 1, HALT!

If read , write 1, HALT!



If read 1, write 0, go right, repeat.

If read 0, write 1, HALT!

If read \square , write 1, HALT!



If read 1, write 0, go right, repeat.

If read 0, write 1, HALT!

If read , write 1, HALT!



If read 1, write 0, go right, repeat.

If read 0, write 1, HALT!

If read , write 1, HALT!



If read 1, write 0, go right, repeat.

If read 0, write 1, HALT!

If read , write 1, HALT!



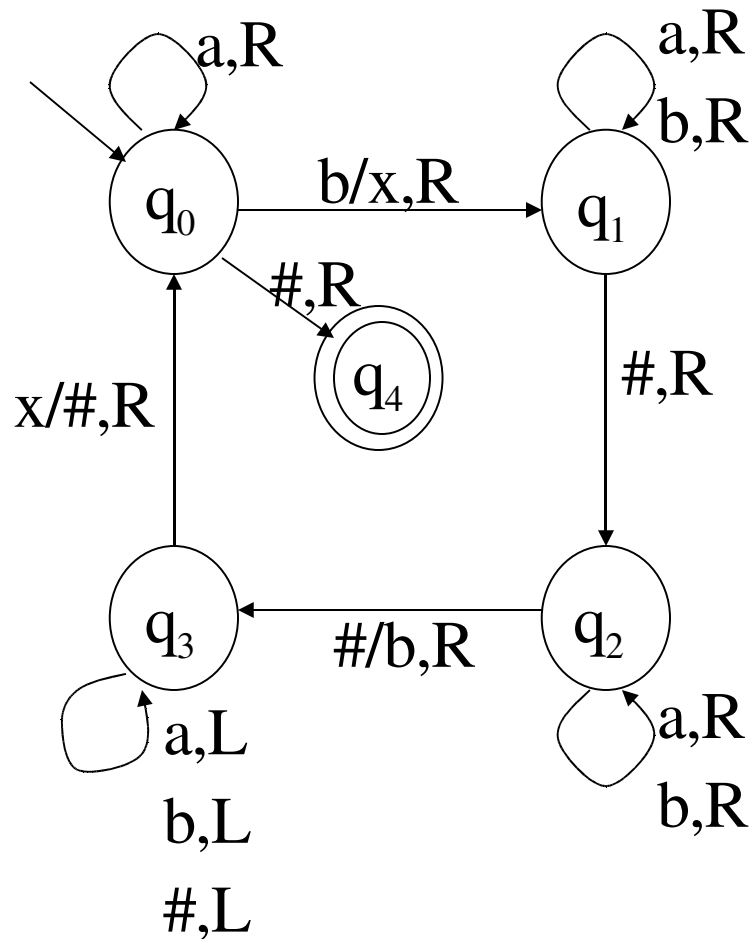
If read 1, write 0, go right, repeat.

If read 0, write 1, HALT!

If read , write 1, HALT!



An Example



$Q = \{q_0, q_1, q_2, q_3, q_4\}$

$\Gamma = \{a, b, x, \#\}$

$B = \#$

$\Sigma = \{a, b\}$

q_0 is the start state

$F = \{q_4\}$

Consider the input:

baaba

What is the output?

Turing Machines

- By using this powerful but simple model, we can start looking at the question of what languages can be defined (equivalently, what problems can be solved) by a computational device. Is there any problem that a computer cannot solve, and what are they?
- We will see in some later classes that there are a lot of them, and they are called “undecidable” problems.

Other Examples

TM can do all sorts of things. Try the followings:

- Add 1 to a unary number. (Easy)
- Add 1 to a binary number.
- Convert a unary number to binary, and vice versa.
- Compare two unary numbers.
- Add two unary numbers.
- Compare two binary numbers x and y . If $x > y$, output 1. Otherwise, output 0.
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TM Simulator

- <http://www.igs.net/~tril/tm/tm.html>
- <http://www.cheransoft.com/vturing/download.html>
- chi scrive un simulatore salta la parte teorica sulle TM